

Optimal Dating Strategy

Wittawat Jitkrittum

Gatsby Tea Talk

22 Oct 2015

Overview



- A.K.A. secretary problem, marriage problem.
- Well known problem in statistics.
- **Goal:** Select the best out of n candidates.

Conditions:

- 1 Candidates arrive one at a time in random order.
- 2 n known.
- 3 Candidates can be ranked unambiguously.
- 4 After the interview (or date), each candidate is either accepted or rejected. Irrevocable decision.
- 5 Decision based only on the relative ranks.

Strategy k

Possibly the only reasonable strategy.

Strategy k (a stopping rule) = “let $k - 1$ go by” strategy

- 1 Sample $k - 1$ candidates to get a rough idea about the pool.
 - 2 Reject them all.
 - 3 Pick the next one who is better than all the $k - 1$ samples.
- **Main question:** What is the best k ?
 - k will depend on n .

What is the best k ?

- k too small. $k = 3$.



- k too big. $k = 14$. Samples include the global best. Last one selected.



- “Success” := finding the global best (after the $k - 1$ samples).
- Let $S_{n,k}$ be the event of success.

$$\text{Best } k = \arg \max_k p(S_{n,k}).$$

- Find k that maximizes the probability of success.

Example: $n = 3$

- Assume $3 > 2 > 1$. Candidate 3 is the best.
- $3! = 6$ possible arrival patterns.
- Green = best candidate selected. Orange = samples.
 - Tables from <http://datagenetics.com/blog/december32012/index.html>

$k = 1$ (no sample)

1	2	3
1	3	2
2	1	3
2	3	1
3	1	2
3	2	1

■ $p(S_{n,k}) = 2/6$

$k = 2$ (1 sample)

1	2	3
1	3	2
2	1	3
2	3	1
3	1	2
3	2	1

■ $p(S_{n,k}) = 3/6$

$k = 3$

1	2	3
1	3	2
2	1	3
2	3	1
3	1	2
3	2	1

■ $p(S_{n,k}) = 2/6$

- $k = 2$ gives the highest probability.

Example: $n = 4$

$k = 1$ (no sample)

1	2	3	4
1	2	4	3
1	3	2	4
1	3	4	2
1	4	2	3
1	4	3	2
2	1	3	4
2	1	4	3
2	3	1	4
2	3	4	1
2	4	1	3
2	4	3	1
3	1	2	4
3	1	4	2
3	2	1	4
3	2	4	1
3	4	1	2
3	4	2	1
4	1	2	3
4	1	3	2
4	2	1	3
4	2	3	1
4	3	1	2
4	3	2	1

$k = 2$

1	2	3	4
1	2	4	3
1	3	2	4
1	3	4	2
1	4	2	3
1	4	3	2
2	1	3	4
2	1	4	3
2	3	1	4
2	3	4	1
2	4	1	3
2	4	3	1
3	1	2	4
3	1	4	2
3	2	1	4
3	2	4	1
3	4	1	2
3	4	2	1
4	1	2	3
4	1	3	2
4	2	1	3
4	2	3	1
4	3	1	2
4	3	2	1

$k = 3$

1	2	3	4
1	2	4	3
1	3	2	4
1	3	4	2
1	4	2	3
1	4	3	2
2	1	3	4
2	1	4	3
2	3	1	4
2	3	4	1
2	4	1	3
2	4	3	1
3	1	2	4
3	1	4	2
3	2	1	4
3	2	4	1
3	4	1	2
3	4	2	1
4	1	2	3
4	1	3	2
4	2	1	3
4	2	3	1
4	3	1	2
4	3	2	1

$k = 4$

1	2	3	4
1	2	4	3
1	3	2	4
1	3	4	2
1	4	2	3
1	4	3	2
2	1	3	4
2	1	4	3
2	3	1	4
2	3	4	1
2	4	1	3
2	4	3	1
3	1	2	4
3	1	4	2
3	2	1	4
3	2	4	1
3	4	1	2
3	4	2	1
4	1	2	3
4	1	3	2
4	2	1	3
4	2	3	1
4	3	1	2
4	3	2	1

$$p(S_{n,k}) = 6/24$$

$$p(S_{n,k}) = 11/24$$

$$p(S_{n,k}) = 10/24$$

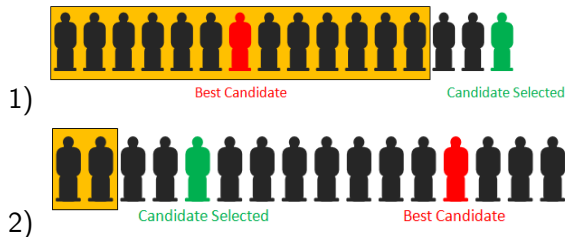
$$p(S_{n,k}) = 6/24$$

■ $k = 2$ gives the highest probability.

■ $k = 1$ or $k = n$ always give uniform probabilities.

Success probability of strategy k

$$\begin{aligned} p(S_{n,k}) &= \sum_{i=1}^n p(i \text{ selected} \cap i \text{ best}) = \sum_{i=1}^n p(i \text{ selected} \mid i \text{ best})p(i \text{ best}) \\ &= \frac{1}{n} \left[\sum_{i=1}^{k-1} \overbrace{p(i \text{ selected} \mid i \text{ best})}^{0. \text{ best is in the samples. See 1)}} + \sum_{i=k}^n \overbrace{p(i \text{ selected} \mid i \text{ best})}^{\text{see 2}} \right] \\ &= \frac{1}{n} \sum_{i=k}^n p(\text{best up to } i-1 \text{ is in the samples} \mid i \text{ best}) \\ &= \frac{1}{n} \sum_{i=k}^n \frac{k-1}{i-1} \end{aligned}$$



As $n \rightarrow \infty$

$$\begin{aligned} p(S_{n,k}) &= \frac{1}{n} \sum_{i=k}^n \frac{k-1}{i-1} = \frac{k-1}{n} \sum_{i=k}^n \frac{n}{i-1} \frac{1}{n} \\ &= \left(\frac{k}{n} - \frac{1}{n} \right) \sum_{i=k}^n \left(\frac{i}{n} - \frac{1}{n} \right)^{-1} \frac{1}{n} \end{aligned}$$

- Let $t := i/n$
- Let $x := k/n =$ proportion of candidates to be used as samples
- $dt := 1/n$
- As $n \rightarrow \infty$, the sum above is a Riemann sum.

$$\begin{aligned} p(S_{n,k}) &= \left(x - \frac{1}{n} \right) \sum_{t=k/n}^{n/n} \left(t - \frac{1}{n} \right)^{-1} \frac{1}{n} \\ &\rightarrow x \int_x^1 \frac{1}{t} dt = -x \log(x). \end{aligned}$$

- $-\frac{\partial x \log x}{\partial x} = 0 \implies x = e^{-1} \approx 0.3679$
- What is the best k ? **Answer:** $k = n/e \approx 37\% \times n$

Practical for dating?

- Say I will live for another 50 years.
- 3 years for a partner. So $n = 50/3 \approx 17$.
- $k = 17/e = 6.25 \approx 6$.
- Spend $3 \times 6 = 18$ years checking the 6 “samples”.
- Then dump them all.
-
- ...
-

Practical for dating?

- Say I will live for another 50 years.
- 3 years for a partner. So $n = 50/3 \approx 17$.
- $k = 17/e = 6.25 \approx 6$.
- Spend $3 \times 6 = 18$ years checking the 6 “samples”.
- Then dump them all.
- Pick the next one that is better than all the 6 people (assuming that I will still be attractive and can be picky).
- ...
-

Practical for dating?

- Say I will live for another 50 years.
- 3 years for a partner. So $n = 50/3 \approx 17$.
- $k = 17/e = 6.25 \approx 6$.
- Spend $3 \times 6 = 18$ years checking the 6 “samples”.
- Then dump them all.
- Pick the next one that is better than all the 6 people (assuming that I will still be attractive and can be picky).
- ... All assuming that I can find the first one to begin with.
-

Practical for dating?

- Say I will live for another 50 years.
- 3 years for a partner. So $n = 50/3 \approx 17$.
- $k = 17/e = 6.25 \approx 6$.
- Spend $3 \times 6 = 18$ years checking the 6 “samples”.
- Then dump them all.
- Pick the next one that is better than all the 6 people (assuming that I will still be attractive and can be picky).
- ... All assuming that I can find the first one to begin with.
- Conclusion: Not practical.

References I

- Took many images from:
<http://datagenetics.com/blog/december32012/index.html>
- Job queue image: <https://www.quora.com/Probability-statistics-1/What-are-the-most-interesting-or-popular-probability-puzzles-in-which-the-intuition-is-contrary-to-the-solution>
- <http://www.math.uah.edu/stat/urn/Secretary.html>
- Thomas S. Ferguson. Statist. Sci. Volume 4, Number 3 (1989), 282-289.